### **Batched Multi-armed Bandits Problem**

Yanjun Han (Stanford EE)

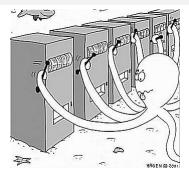
Joint work with:

Zijun Gao Zhimei Ren Zhengqing Zhou Stanford Stats Stanford Stats Stanford Math

NeurIPS 2019, Vancouver, Canada

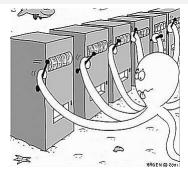
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- sequential decision making
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- action space: K arms
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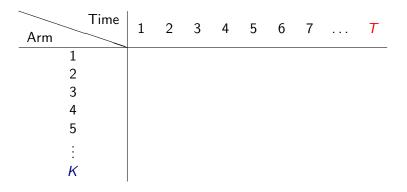
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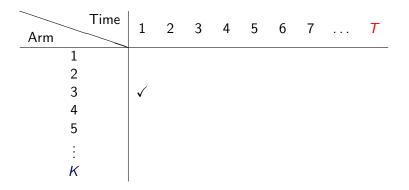


#### Space Domain: Bandit Feedback

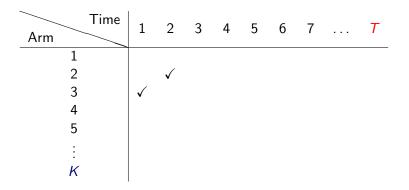
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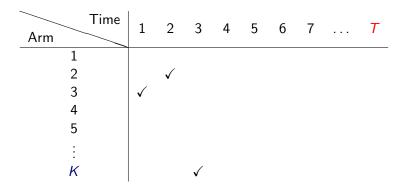
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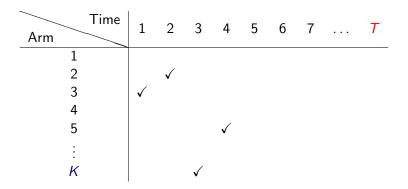
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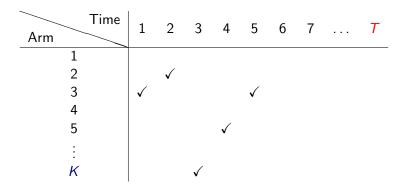
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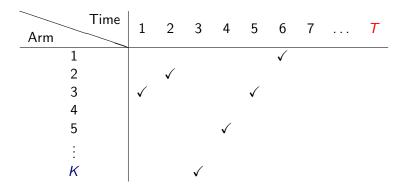
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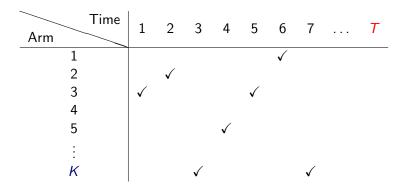
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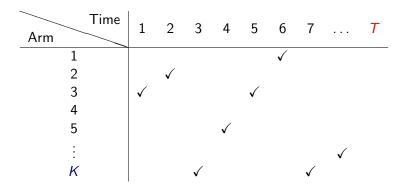
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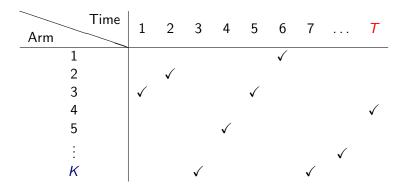
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Clinical trial



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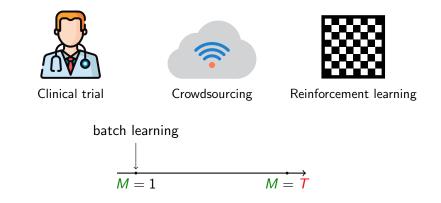
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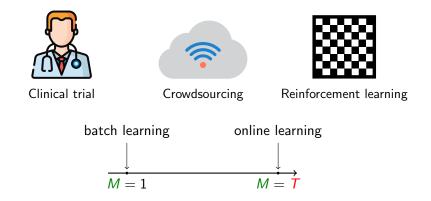
Reinforcement learning

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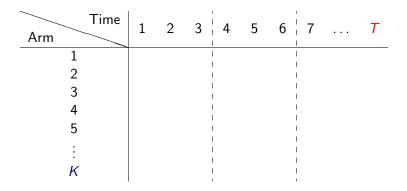


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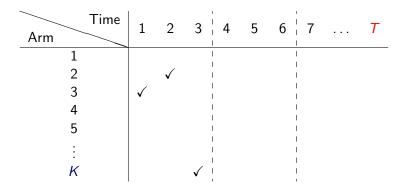


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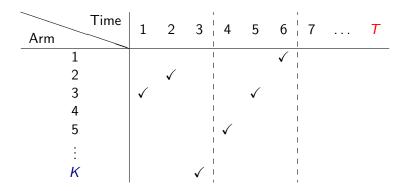
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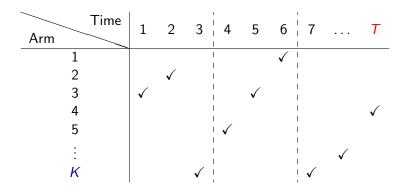
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- adaptive grid: the next grid point determined by historic data
- task: design policy + grid

## Two Types of Regrets

Tight analysis of stochastic MAB [Vog'60, LR'85, AB'09]:

$$\mathbb{E}[R(\pi^1)] \leq C \cdot \sqrt{K T} \ \mathbb{E}[R(\pi^2)] \leq C \cdot \sum_{i 
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#### Problem-dependent Regret

$$R_{\text{pro-dep}}(K, M, \mathbf{T}) = \inf_{\pi, \mathcal{T}} \sup_{\Delta > 0} \Delta \cdot \sup_{\Delta_i \in \{0\} \cup [\Delta, \sqrt{K}]} \mathbb{E}[R(\pi)]$$

Full online case:

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Required number of batches [ACBF'02, CBDS'13]:

$$\begin{aligned} R_{\min-\max}(K, \log T, T) &= \widetilde{\Theta}(\sqrt{KT}) \quad (\text{UCB2}) \\ R_{\min-\max}(K, \log \log T, T) &= \widetilde{\Theta}(\sqrt{KT}) \quad (\text{switching cost}) \end{aligned}$$

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Two-armed case with static grid [PRCS'16]:

$$R_{\text{min-max}}(2, M, T) = \widetilde{\Theta}(T^{\frac{1}{2-2^{1-M}}})$$
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Lower bounds typically very challenging [JJNZ'16, AAAK'17, DRY'18, ...].

### Theorem 1 (Upper Bound)

There exist policies  $\pi^1,\pi^2$  such that

$$\mathbb{E}[R(\pi^{1})] \leq \mathsf{polylog}(K, T) \cdot \sqrt{KT^{\frac{1}{2-2^{1-K}}}}$$
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M = log log T batches sufficient for centralized minimax regret
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### Minimax Grid

 $\mathcal{T}_{\mathsf{minimax}} = \{t_1, \cdots, t_M\}$  with

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### Main Result II: Static Lower Bound

Theorem 2 (Static Lower Bound)

Under any static grid,

$$R_{\text{min-max}}(K, M, T) = \Omega(\sqrt{KT}^{\frac{1}{2-2^{1-M}}})$$
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- match the upper bounds within logarithmic factors
- proof uses a max-min approach: find multiple fixed reward distributions under which no policy performs uniformly well

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#### Indistinguishability Lemma

Let  $Q_1, \dots, Q_n$  be probability measures on some common probability space. Then for any tree T = ([n], E) and test  $\Psi$ ,

$$\frac{1}{n}\sum_{i=1}^n Q_i(\Psi \neq i) \geq \sum_{(i,j)\in E} \frac{1}{2n} \exp(-D_{\mathrm{KL}}(Q_i || Q_j)).$$

Theorem 3 (Adaptive Lower Bound)

Under any adaptive grid,

$$R_{\text{min-max}}(K, M, T) = \Omega(M^{-2} \cdot \sqrt{KT^{\frac{1}{2-2^{1-M}}}})$$
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- use a min-max approach instead: construct corresponding reward distributions after a policy is given

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Construct reward distributions  $P_1, P_2, \cdots, P_M$  and events  $A_1, \cdots, A_M$ .

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### Lemma 2 (Covering of Events)

For any policy it holds that

$$\sum_{m=1}^M P_m(A_m) \geq \frac{1}{2}.$$

Take-home message:

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### Thank you!